

# PROOF OF ORE'S THEOREM

CHING-HAO, WANG

## 1. ORE'S THEOREM

In this note, we prove the following result:

**Theorem 1.** *Suppose that  $G$  is a simple graph with  $n$  vertices,  $n \geq 3$  and  $\deg(x) + \deg(y) \geq n$  whenever  $x$  and  $y$  are non-adjacent vertices in  $G$ , then  $G$  has a Hamilton circuit.*

*Proof.* We prove by contradiction:

(a) Suppose that  $\deg(x) + \deg(y) \geq n$  for every pair of nonadjacent vertices  $x$  and  $y$  in  $G$ . If  $G$  does not have a Hamilton circuit, continue as long as possible in adding missing edges to  $G$  (since  $G$  is not necessarily a complete graph) one at a time in such a way that we do not obtain a graph with a Hamilton circuit. This cannot go on forever, because once we've formed the complete graph by adding all missing edges, there's a Hamilton circuit. Whenever the process stops, we have obtained a (necessarily non-complete) graph  $H$  with the desired property.

(b) If we add one more edge to  $H$ , this produces a Hamilton circuit, which uses all added edges. The path consisting of this circuit with the added edge omitted is the Hamilton path in  $H$ .

(c) Let  $v_1, v_2, \dots, v_n$  be a Hamilton path in  $H$ . Clearly  $v_1$  and  $v_n$  are not adjacent in  $H$ , because  $H$  has no Hamilton circuit. Therefore they are not adjacent in  $G$ . By hypothesis,

$$\deg(v_1) + \deg(v_n) \geq n$$

We rewrite in the other form,

$$(n-1) - \deg(v_n) \leq \deg(v_1) - 1$$

, the left-hand-side (LHS) is just the number of vertices not adjacent to  $v_n$  (not include  $v_n$  itself).

(d) Let  $S$  be the set of vertices preceding each vertex adjacent to  $v_1$  in the Hamilton path. Because there's no vertex following  $v_n$ , so  $v_n \notin S$ . Each one of the  $\deg(v_1)$  vertices adjacent to  $v_1$  gives rise to one element in  $S$ . So  $|S| = \deg(v_1)$

(e) By (c), there are at most  $\deg(v_1) - 1$  vertices other than  $v_n$  not adjacent to  $v_n$ , and by (d) there are  $\deg(v_1)$  vertices in  $S$ , none of which is  $v_n$ . Therefore at least one vertex of  $S$  is adjacent to  $v_n$  (since each of the  $\deg(v_1)$  vertices in  $S$  is either adjacent or not adjacent to  $v_n$ , but there are at most  $\deg(v_1) - 1$  vertices

not adjacent to  $v_n$ ). Let  $v_k$  be such a vertex and, by definition,  $H$  contains the edges  $\{v_k v_n\}$ ,  $\{v_{k+1} v_1\}$ , where  $1 < k < n_1$ .

(f) Now  $v_1, v_2, \dots, v_{k-1}, v_k, v_n, v_{n-1}, \dots, v_{k+1}, v_1$  is a Hamilton circuit in  $H$ , which contradicts our construction of  $H$ . Therefore our assumption that  $G$  did not originally have a Hamilton circuit is wrong, and our proof by contradiction is complete.  $\square$

#### REFERENCES

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DEPARTMENT OF ELECTRICAL ENGINEERING, NATIONAL TSING-HUA UNIVERSITY, HSINCHU, TAIWAN, R.O.C., 30013